Ramsey Games: Avoiding Triangles in Graphs

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University of Minnesota - Twin Cities

MathPath 2013

Dr. B. (UMN)

Ramsey Games

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Alexander Ramsey, first Governor of the Minnesota Territory

The Wrong Ramsey



Alexander Ramsey, first Governor of the Minnesota Territory

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The Wrong Ramsey



Alexander Ramsey, first Governor of the Minnesota Territory The Right Ramsey



Frank Plumpton Ramsey

"How big does something have to be for some thing I'm looking for to show up inside it?"

The Right Ramsey



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<u>Not-much-fun fact:</u> Ramsey died at the age of 26 from complications after an abdominal operation to treat his jaundice.

The Right Ramsey



Frank Plumpton Ramsey

What's the Pattern?

1, 2, 31, 4, 7, 10, 132, 6, 10, 14 What's the Pattern?

1, 2, 31, 4, 7, 10, 132, 6, 10, 14

These are called Arithmetic Progressions.

What's the Pattern?

1, 2, 3 1, 4, 7, 10, 13

2, 6, 10, 14

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Just to be formal:

Definition

An Arithmetic Progression of length k (called k-AP for short) is a sequence of k numbers such that the difference between any consecutive numbers is constant.

Avoiding Arithmetic Progressions

... to the document camera!

Dr. B. (UMN)

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In 1927, proved "Van der Waerden's Theorem"

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What has he got to do with Arithmetic Progressions?

Theorem

For any natural number k there exists some number n large enough that any r-coloring of the integers 1, 2, ..., n must contain either a red k-AP or a blue k-AP.

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For k = 3 and r = 2, it turns out that 9 is big enough.

Van der Waerden Numbers

Just for fun...

	3	4	5	6	7
2 colors	9	35	178	$1,\!132$	
3 colors	27	293			
4 colors	76				
5 colors					
6 colors					

Van der Waerden Numbers

Just for fun...

	3	4	5	6	7
2 colors	9	35	178	$1,\!132$	> 3,703
3 colors	27	293	> 2,173	> 11, 191	> 48,811
4 colors	76	> 1,048	> 17,705	> 91,331	> 420, 217
5 colors	> 170	> 2,254	> 98,740	> 540,035	
6 colors	> 223	> 9,778	> 98,748	> 816,981	

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Lunchtime challenge: can you 3-color the numbers $1, 2, \ldots, 26$ without making a monochromatic 3-AP?

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1, 2, 3

1, 3, 5, 7

Dinnertime challenge: what is the longest arithmetic progression of prime numbers you can find?

Here is the longest one currently known:

 $43, 142, 746, 595, 714, 191 + 23, 681, 770 \cdot 223, 092, 870 \cdot n,$

for n = 0 to 25.

Where are the graph games?

So much for Arithmetic Progressions.

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So much for Arithmetic Progressions.

What happens when a mathematician reads chemistry papers?



James Joseph Sylvester

Dr. B. (UMN)



James Joseph Sylvester





James Joseph Sylvester

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- The degree, d(v), of a vertex is the number of edges containing it.
- The max degree of G, $\Delta(G)$, is $\max\{d(v) : v \in V(G)\}.$

Some important graphs



Some important graphs



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- The complete graph, K_{ℓ} , has ℓ vertices, and every two vertices form an edge.
- The cycle, C_{ℓ} , has vertices v_1, \ldots, v_{ℓ} , and $v_i v_j$ is an edge iff $|i j| = 1 \pmod{\ell}$.

Ramsey + Sylvester = \dots

Graph Ramsey Theory.

Challenges: (back to the document camera)

Let's color edges instead of vertices this time.

• Game 1: I will give you a graph and you have to color the edges to avoid a triangle.

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Which game would you rather play? Why?

Game 1

You should not play this game! Ramsey says:

Theorem

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Theorem

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In fact, a similar statement is true no matter what graph we are trying to avoid and how many colors we get to use. • Game 2: Same as Game 1, but the graph's maximum degree will be at most 4.

Game 2 vs Game 3 $\,$

- Game 2: Same as Game 1, but the graph's maximum degree will be at most 4.
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I have a lot more power in Game 3! Let's play... (to the document camera!)

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R(6) is somewhere between 102 and 165 We would have better luck trying to destroy the aliens!

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Many variations

• What if we want to avoid one graph in red and a different graph in blue?

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- What if I promise to stay in some other family of graphs (instead of max-degree-4 graphs)?

- What if we want to avoid one graph in red and a different graph in blue?
- What if I promise to stay in some other family of graphs (instead of max-degree-4 graphs)?
- What if I am replaced by a random opponent?

Do you want to know more?

Breakout: Graph Ramsey Games, Week 4